Problem 1. ★

We want to implement a parametric function reverse#(n) that reverses the bits of its n-bit input argument. For example, if Bit#(4) x = {a, b, c, d}, then reverse#(4)(x) should return {d, c, b, a}.

(A) Implement reverse#(n) by recursing on the parameter n (i.e., calling reverse#(k) with k < n). You cannot use a for loop.

```
function Bit#(n) reverse#(Integer n)(Bit#(n) x);

endfunction
```

(B) Implement reverse#(n) using a for loop.

```
function Bit#(n) reverse#(Integer n)(Bit#(n) x);

endfunction
```
Problem 2.

Parameterize the bit-scan-reverse function from Lab 3 to take as input an n-bit vector and output the index of the first non-zero bit scanned from the largest index (i.e., the position of the most-significant 1). Assume that the parameter n is a power of 2.

(A) Implement `bitScanReverse#(n)` without using a `for` loop.

```verilog
function Bit#(log2(n)) bitScanReverse#(Integer n)(Bit#(n) x);
endfunction
```

(B) Implement `bitScanReverse#(n)` using a `for` loop.

```verilog
function Bit#(log2(n)) bitScanReverse#(Integer n)(Bit#(n) x);
endfunction
```

(C) When synthesized manually (i.e., without logic optimizations, just the gates that your circuit expresses), how does propagation delay grow with the number of input bits for each implementation? Use order-of notation.
Problem 3. ★

In Lab 3, we wrote a function `isPowerOfTwo` that computes whether a 4-bit input was a power of 2 or not. This function checks whether there is only one bit in the input that is equal to 1.

We want to generalize `isPowerOfTwo` by rewriting it as a parametric function that works with inputs of arbitrary bit-width.

(A) Implement `isPowerOfTwo#(n)` using a `for` loop. Do not use addition to count up the bits of the input, which would be inefficient.

```verilog
function Bool isPowerOfTwo#(Integer n)(Bit#(n) x);
endfunction
```

(B) If your implementation above has $\Theta(n)$ propagation delay when synthesized manually (i.e., without logic optimizations, just the gates that your circuit expresses), then rewrite `isPowerOfTwo#(n)` so that it has $\Theta(\log n)$ propagation delay.

*Hint:* Since you used a `for` loop above, you probably have a linear chain of gates in your design. Instead, think about how to solve this problem by composing functions so that, at each step, you halve the number of input bits each function processes. This will produce a tree of gates with logarithmic depth. You’ll likely need to use an auxiliary parametric function that recurses on its own parameter.
Problem 4. (adapted from Quiz 1 Fall 2018, 20/100 points)
(NOTE: Part C of this question will be much easier once you have completed Lab 4)

(A) (5 points) The following parametric function \( f \) performs a basic operation using \( a \) and \( b \). We want \( f_2 \) to implement the same function as \( f \). Fill in the blank in \( f_2 \) to make the two functions equivalent. Write a single-line expression that uses the ternary operator (?) :.

\[
\text{function} \quad \text{Bit#}(n) \quad f#(\text{Integer } n)(\text{Bit#}(n) \ a, \ \text{Bit#}(1) \ b);
\quad \text{Bit#}(n) \ x = 0;
\quad \text{for} \text{ (Integer } i = 0 ; i < n ; i = i+1) \text{ begin}
\quad \quad x[i] = a[i] \ ^ \ b;
\quad \end{endfunction}

\[
\text{function} \quad \text{Bit#}(n) \quad f2#(\text{Integer } n)(\text{Bit#}(n) \ a, \ \text{Bit#}(1) \ b);
\quad \text{return} \ (\quad \quad ) \ ? \quad \quad : \quad \quad ;
\endfunction
\]

(B) (5 points) Write the truth table for the combinational device described by the function below.

\[
\text{function} \quad \text{Bit#}(2) \quad f(\text{Bit#}(1) \ a, \ \text{Bit#}(1) \ b, \ \text{Bit#}(1) \ c);
\quad \text{Bit#}(2) \ ret = \text{zeroExtend}(a) + \text{signExtend}(b);
\quad \text{case} \ \{(a,b)\}
\quad \quad 0: \ ret = \{1, c\};
\quad \quad 2: \ ret = \{a \ ^ \ b, a \ & \ b\};
\quad \quad 3: \ ret = \sim \text{signExtend}(c);
\quad \endcase
\quad \text{return} \ ret;
\endfunction
\]

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<tr>
<th>a</th>
<th>b</th>
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(C) (5 points) The following parametric function \( g \) performs a specific arithmetic operation on \( n \)-bit operands \( a \) and \( b \). We want the function \( g2 \) to implement \( g \) in a single line of code. Fill in the blank with a single expression to make \( g2 \) equivalent to \( g \).

```haskell
function Bit#(1) g#(Integer n)(Bit#(n) a, Bit#(n) b);
    Bit#(2) ret = 'b10;
    for (Integer i = n-1 ; i >= 0 ; i = i-1) begin
        if ({a[i], b[i]} == 'b01) ret = {0, ret[1] | ret[0]};
        else if ({a[i], b[i]} == 'b10) ret = {0, ret[0]};
    end
    return ret[1] | ret[0];
endfunction

function Bit#(1) g2#(Integer n)(Bit#(n) a, Bit#(n) b);
    return ________________________________;
endfunction
```

(D) (5 points) Finish the following circuit diagram to implement function `computeB`, given below. You may only use 32-bit 2-to-1 multiplexers, constants (0, 1, 2, 3, ...) and logic gates (AND, NOT, OR, XOR). We have provided three 32-bit greater-than-or-equal (>=) comparators for you.

```haskell
function Bit#(32) computeB(Bit#(32) in);
    Bit#(32) out = 0;
    if ( in >= 1 ) out = 1;
    if ( in >= 5 ) out = 5;
    if ( in >= 10 ) out = 10;
    return out;
endfunction
```

![Circuit Diagram](image)